

- A general multi-affine system in \mathbb{R}^S ($S = \# \text{ species}$) can have multiple equilibria.

- These sys's contain only constant, linear, and bilinear terms.

- (ex) 2D system ($S=2$): a multi-affine vector field is

$$\dot{x}(x_1, x_2) = c_1 + c_2 x_1 + c_3 x_2 + c_4 x_1 x_2$$

Chemical Reaction Network Theory

[Martin Feinberg, Chemical Engineering Science 42(10), 1987
and Archive for Rational Mechanics and Analysis, vol.
132, 1995]

- Provides techniques for determining the existence, multiplicity, and stability properties of equilibria of nonlinear dynamical systems of certain classes.
- Results apply to networks with high-dimensional, complex dynamics; sometimes hold regardless of values of system parameters.
- "Deficiency theorems" state results for equilibria + cyclic trajectories in the positive orthant of state space. Results apply to closed, spatially homogeneous systems that have constant volume; may be extended to open systems by including "pseudo reactions" of the form $O \rightarrow X_i, X_i \rightarrow O$.

(2)

Linear Macroscopic Model

$$\dot{\underline{x}} = -\underline{K} \underline{x} \quad \text{conservation constraint: } \underline{1}^T \underline{x} = 1$$

- CRN contains reactions only of type $x_i \xrightarrow{k_{ij}} x_j$
- Models scenario where robots reallocate themselves among a set of tasks that are to be executed in parallel, continuously, and independently of one another.

$x_i(t)$ = population fraction of robots performing task i at time t (a continuous quantity)

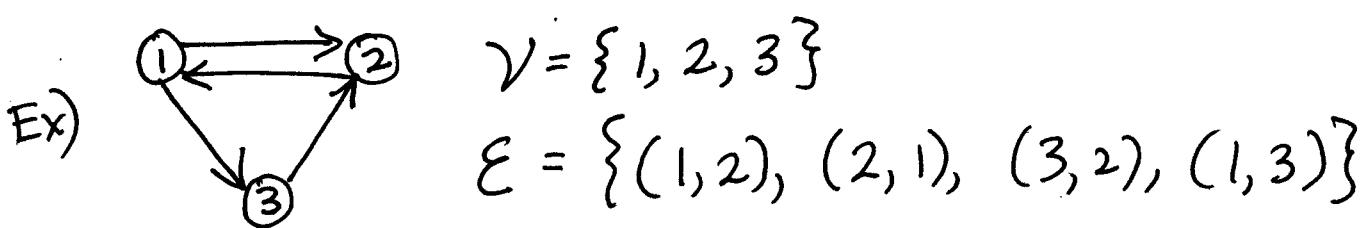
$$\underline{x} \in \mathbb{R}^S \quad S = \# \text{ of tasks}$$

$$\underline{x} = [x_1 \ x_2 \ \dots \ x_{S-1} \ x_S]^T$$

Vertex set $\mathcal{V} = \{1, 2, \dots, S-1, S\}$ = set of tasks

Edge set \mathcal{E} = set of possible transitions between tasks

- If a robot can switch from task i to task j , then we say that tasks i and j are adjacent, denoted by $i \sim j$ or by (i, j)
- \mathcal{E} is the set of all (i, j) task pairs



(3)

Graph $G = \{V, E\}$ models precedence constraints between tasks.

- If there is a sequence of vertices $\{v_0, v_1, \dots, v_p\} \in V$ such that $v_0 = i$, $v_p = j$, and $(v_{k-1}, v_k) \in E$, $k = 1, \dots, p$ for every $\{i, j\} \in V$, then G is strongly connected. ($\{v_0, v_1, \dots, v_p\}$ is called a directed path.)
- The example graph is strongly connected.

We will look at some properties of strongly connected graphs that are useful for task allocation in swarms.

Theorem: If the graph G is strongly connected, then the model $\dot{\underline{x}} = -\underline{K}\underline{x}$ subject to $\underline{1}^T \underline{x} = 1$ has a unique stable equilibrium.

Proof: The equilibrium \underline{x}^e is the solution of

$$\underline{K}\underline{x}^e = \underline{0} \text{ subject to } \underline{1}^T \underline{x}^e = 1.$$

\underline{x}^e eigenvector belonging to the zero eigenvalue

- G is strongly connected \Rightarrow rank of $\underline{K} = S-1$

$$\underline{K} \in \mathbb{R}^{S \times S} \Rightarrow \text{rank}(\underline{K}) + \text{nullity}(\underline{K}) = S$$

$$\Rightarrow S-1 + \text{nullity}(\underline{K}) = S$$

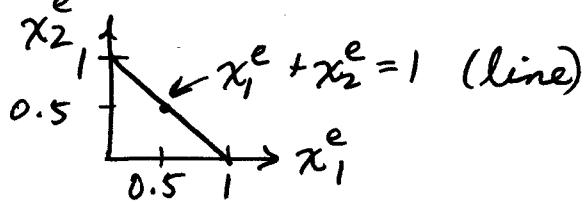
$$\Rightarrow \text{nullity}(\underline{K}) = 1 \Rightarrow \text{null space of } \underline{K} \text{ has dimension 1}$$

$\underline{x}^e \in \text{Null space of } \underline{K} \Rightarrow \underline{x}^e$ can take values along a line in \mathbb{R}^S

(4)

The conservation constraint $\underline{1}^T \underline{x}^e = 1$ defines an (≤ 1) -dimensional hyperplane in \mathbb{R}^S .

ex) $S=2 : \underline{1}^T \underline{x}^e = 1 \Rightarrow x_1^e + x_2^e = 1$



\Rightarrow the null space of \underline{K} must intersect this hyperplane.

Null space is 1D \Rightarrow intersects plane at a point.

$\Rightarrow \underline{x}^e$ can have a single (unique) value, which occurs at this point.

\Rightarrow the model has a unique equilibrium point.

Now consider $\underline{I} = t \underline{I} - \underline{K}$, where $t > 0$ and

$\underline{I} \in \mathbb{R}^{S \times S}$ is the identity matrix.

Assume that t is large enough such that \underline{I} is a nonnegative matrix (all its elements are ≥ 0).

- Since G is strongly connected, the matrix $-\underline{K}$, and therefore \underline{I} , is irreducible. via a permutation
 - An irreducible matrix is not similar¹ to a block upper triangular matrix, i.e.

$$\begin{bmatrix} \underline{K}_1 & | & \underline{K}_2 \\ \hdashline \text{---} & | & \text{---} \\ \underline{0} & | & \underline{K}_3 \end{bmatrix}$$

- \underline{K} is irreducible if and only if G is strongly connected.]

(5)

\underline{I} is nonnegative and irreducible

\Rightarrow By the Perron-Frobenius theorem, \underline{I} has a real, positive, simple eigenvalue $\lambda_m(\underline{I})$ such that all other eigenvalues of \underline{I} , $\lambda(\underline{I})$, satisfy $|\lambda(\underline{I})| < \lambda_m(\underline{I})$.

This eigenvalue also satisfies :

$$\min_j \sum_{i=1}^s I_{ij} \leq \lambda_m(\underline{I}) \leq \max_j \sum_{i=1}^s I_{ij}$$

Columns of \underline{K} sum to 0 $\Rightarrow \sum_{i=1}^s I_{ij} = t - \underbrace{\sum_{i=1}^s K_{ij}}_{=0}$
 \Rightarrow both sides of the inequalities = t

$$\Rightarrow t \leq \lambda_m(\underline{I}) \leq t \Rightarrow \lambda_m(\underline{I}) = t$$

• $\lambda(\underline{I}) = \lambda(-\underline{K}) + t$

$$\Rightarrow t = \lambda_m(-\underline{K}) + t \Rightarrow \lambda_m(-\underline{K}) = 0$$

$$|\lambda(\underline{I})| = |\underbrace{\lambda(-\underline{K})}_{} + t| < \lambda_m(\underline{I}) = t$$

\Rightarrow all besides $\lambda_m(-\underline{K})$ have $\operatorname{Re}(\lambda(-\underline{K})) < 0$.

\Rightarrow Eigenvalue of $-\underline{K}$ corresponding to $\lambda_m(\underline{I})$ is 0,
 All other eigenvalues of $-\underline{K}$ have negative real parts. \quad a simple eigenvalue

$\Rightarrow \underline{x}^e$ is a stable equilibrium point.

• non-zero off-diag entries of $-\underline{K}$ are $> 0 \Rightarrow$ there is a solution $\underline{x} > \underline{0}$.